

## How **You** in HEP can benefit from Functional Programming

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## Contents

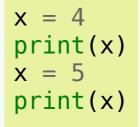
- What is **Functional Programming** (FP)
- Relation between the FP and High Energy Physics communities
- FP in C++ and Python for your everyday work
- FP-inspired techniques in data analysis
- Some easy FP paradigms to **improve your code**

## Introduction

- Who of you has heard about Functional Programming before?
- Tell me what you think of this Python code:
- Tell me what you think of this C++ code:

```
int n = 10;
for(int i = 0; i < n; ++i)
{
    std::cout << i << " ";
}
std::cout << std::endl;</pre>
```

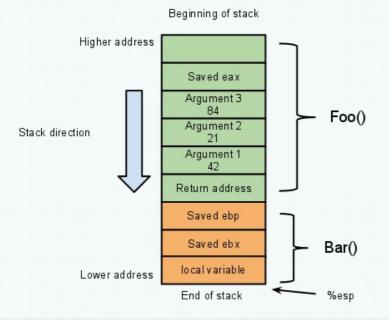
• Lies, lies lies! You can't trust the "definition" of variables!

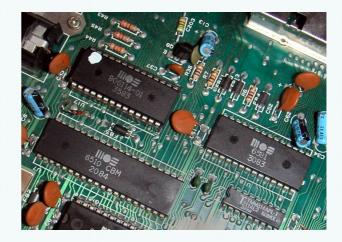




## But why do we lie so much?

- A historical sketch of programming:
- From machine code to assembly
- From assembly to e.g. Fortran or C

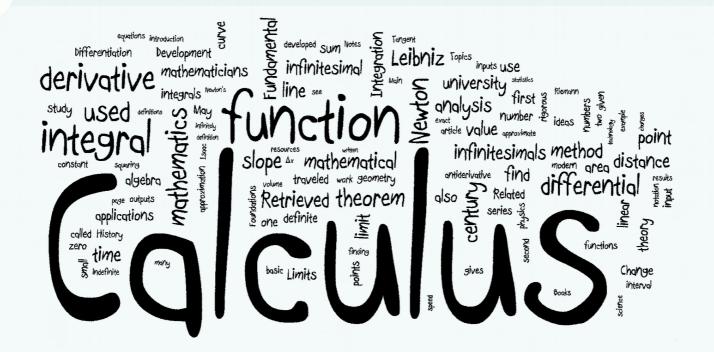




- We got used to variables standing for a location in memory...
- ...not as short-hands for values as in mathematics!
- What is the next level of abstraction?



## **Functional Programming**



- Treat computation as the evaluation of mathematical functions
- Avoid changing-state and mutable data
- Higher level of abstraction requires very smart compilers



## The Lambda Calculus

- Lambda calculus: a formal system for expression computation based on functions as a universal model of computation
- Very interesting research topic that plays an important role in FP

Syntax	Name	Description
X	Variable	A character or string representing a parameter or mathematical/logical value
(λx.M)	Abstraction	Function definition (M is a lambda term). The variable x becomes bound in the expression.
(M N)	Application	Applying a function to an argument. M and N are lambda terms.
Operation	Name	Description
$(\lambda x.M[x]) \rightarrow (\lambda y.M[y])$	α-conversion	Renaming the bound (formal) variables in the expression. Used to avoid name collisions.
$((\lambda x.M) E) \rightarrow (M[x:=E])$	β-reduction	Replacing the bound variable with the argument expression in the body of the abstraction

## Why was I showing this?

- Consider trivial expression  $(\lambda x.x)$
- In Python it would be lambda x: x
- That should look familiar! Lambda functions are now part of many languages so you can create unnamed "throw-away" functions
- Example: map(lambda x: x\*\*2, [1,2,3])
- We just saw another FP concept: Higher-order functions!
   map is a function that takes a function as an argument

Now you know why it's called "lambda function" in Python. But why "lambda calculus in the first place?

• Recursive problem... which is a good keyword



## **Recursion in Lambda Calculus**

#### How to implement recursion?

- Not possible to refer to the definition of a function in a function body!
- What we need is some higher-order function to apply on a function to call it recursively:

fix f = f(fix f), and therefore  $fix f = f(f(\dots f(fix f) \dots))$ .

 This fixed-point combinator can be implemented in lambda Calculus by Haskell Curry's Y combinator:

 $Y = \lambda f.(\lambda x.f(x x))(\lambda x.f(x x))$ 

## Why was I showing this again?

- The US company **Y Combinator** is arguably the most famous startup incubator in silicon valley
- It invested in Dropbox, Airbnb, Reddit, Docker, etc.
- If FP concepts inspire trendy company names now, we are definitely up to something!

## **Y** Combinator



## **Functional Programming Languages**

• The most famous one is **Haskell**, which sticks close to the **typed lambda calculus**. These are the strong points of the language:

#### - Purity:

functions have no side-effects and outputs only depends on inputs

#### - Strong typing:

all variable types must be known at compile-time and no implicit casting, which results in most bugs being caught at compile time and enables powerful static code analysis

#### - Elegance:

code often reads like a high-level algorithm description

#### – Laziness:

nothing is evaluated until it has to be evaluated

- In particular the first two bullet-points result in fewer bugs than you'd normally get in e.g. C++ and Python
- I will not go deeper in the topic of FP languages and will explain how to use these concepts in C++ or Python instead

## Functional Programming in the Real World

- Some companies use pure FP languages like **Haskell** or **Ocaml** (the latter developed by **INRIA**)
- Mostly in silicon valley, startup scene, fin-tech sector and France





- Companies like FP languages for:
  - Robustness
  - Scalability
  - Rapid development times (once you get the hang of it)
  - Readability of code
  - It attracts high quality nerds

## Functional Programming in High Energy Physics

#### Google searches:

- physics+"object oriented programming"
   2 Million results
- physics+"functional programing" 7000 results
- Conclusion: we definitely didn't jump on that bandwagon...
- High energy physics has an important Fortran legacy
  - Fortran: as un-functional as it gets. Variables all have to be declared in the beginning and then you mutate them heavily!
  - Many people still write code like this today
- Fortunately, we use a lot of C++ and Python, which now took over many concepts from PF for example:
  - Powerful type systems (C++)
  - Lambda functions and higher-order functions (both)
  - Mutable variables can be avoided without loosing much speed thanks to powerful optimizers in moders compilers (C++)



# What you can take from FP to write better code **Today**

- Stick to pure functions as much as possible.
  - A pure functions output only depends on the input (not on some external state) and has no "side effects" (implicit effects on external state)
- Ensure that your variables have well-defined types
  - In Python you can't really do that
  - In C++ that means avoid nullptr , casts, void\* and violations of the Liskov Substitution principle
- Try to find elegant, abstract solutions to your problems with powerful higher-level functions and define your own
  - In C++ that means get familiar with function pointers
  - In Python that means use libraries with a high level of functional abstraction like Pandas or sklearn
- Don't be afraid of avoiding mutable data
  - Don't try to mutate variables to "make your program more efficient"
  - Infest your C++ code with const

FUNCTION f:

OUTPUT f(x)

## Making a case for **Pure Functions**

- Typical C++ HEP coding pattern on the right
- Often huge "god classes" which have a lot of member variables that get mutated all over the place
- That's dangerous because information may "leak over" to other event because you didn't reset a variable for example
- You can never be sure what the actual inputs and outputs of a function are without reading it's body
- Don't be that guy

```
class MyAnalyzer() {
```

```
public:
```

MyAnalyzer(Configuration config);

run(Event iEvent); // for event loop

private:

};

// Some functions that do

- // something to the class members.
- // Real inputs and outputs NOT CLEAR!

```
void selectElectrons(float cut);
void makeSuperClusters(bool keepBrems);
bool fitTrack(int nHits);
```

```
// Metric shit-ton of variables
// which all get mutated in
// the class member functions
// (hidded side effects...)
int iEvent_;
float pt_;
float eta_;
float phi_;
// ...
// ...
```

## Making a case for Type Safety

• Very nice negative example is the ROOT TTree's **Branch** method:

TBranch \* TTree::Branch ( const char \* name, void \* address, const char \* leaflist, Int\_t bufsize = 32000 );

 Problem is the void pointer. Type of address is instead specified as argument, e.g.:

> Float\_t mass; tree→Branch("mass", &mass, "mass/F");

- What if you later change the type of mass in the code from "Float\_t" to "Double\_t" and forget to change "F" to "D"? The compiler won't complain and you'll get garbage in your tree
- Why not using the compiler and type system as your ally to detect bugs early? Missed opportunity here!

```
#include "TFile.h"
#include "TRandom.h"
#include "TFile.h"
#include "Ttree.h"
#include <vector>
#include <cmath>
class MySimulator {
  public:
    MySimulator() {
        tree = new TTree("tree", "tree");
        tree ->Branch("n"
                              , &nElectrons , "n/i");
        tree ->Branch("eta"
                              , &etaVec );
        tree_->Branch("phi" , &phiVec_);
        tree_->Branch("pt"
                              , &ptVec );
        rng = new TRandom(); // random number generator
    }
    ~MySimulator() {
        tree ->Write();
        delete rng ;
    }
    void simulate(int iEvent) {
        etaVec .clear(); // don't forget to clear vectors at each event
        phiVec .clear();
        ptVec .clear();
        nElectrons_ = 4 + rng_->Integer(4); // how many electrons?
        for(int iElectron = 0; iElectron < nElectrons ; ++iElectron) {</pre>
            phiVec_.push_back( rng_->Uniform() * 2.*M_PI - M_PI );
            etaVec .push back( rng ->Uniform() * 5.0 - 2.5
            ptVec .push back( rng \rightarrow Exp(10.) + 20.
                                                                  );
        }
        tree ->Fill(); // fill the TTree with event information
    }
  private:
    TTree* tree ;
    TRandom* rng ;
    unsigned int nElectrons ; // variables for TTree branches
    std::vector<float> etaVec ;
    std::vector<float> phiVec ;
    std::vector<float> ptVec ;
};
int main() {
    int nEvents = 10000; // number of events to simulate
    TFile* file = TFile::Open("data.root", "RECREATE");
    MySimulator simulator{};
    for(int iEvent = 0; iEvent < nEvents; ++iEvent) {</pre>
        simulator.simulate(iEvent);
```

);

## C++ case study: filling a TTree

- Let's try to find the "non-FP" • parts together:
  - One mutated variable per tree \_ branch
  - Impure TTree::Fill() and TTree::Write() functions
  - Mutated TRandom
  - Vectors mutated all over the place:
    - std::vector<T>::clear()
    - std::vector<T>::push back()



delete file:

```
#ifndef __S_TREE__
#define S TREE
#include "TTree.h"
#include <mutex>
template<class T>
void addTreeBranch(TTree* tree, const char* name, T& variable) {
    tree->Branch(name, &variable);
#define TYPES WITH SUFFIX \
                 , "/B") X(unsigned char , "/b") \
    , "/S") X(unsigned short, "/s") \
    , "/I") X(unsigned int , "/i") \
     X(char
     X(short
     X(int
                    , "/L") X(unsigned long , "/l") \
     X(long
                      , "/F") X(double
     X(float
                                                , "/D") ∖
     X(char *
                      , "/C") X(bool , "/O")
#define X(Type, suffix) \
template<> \
void addTreeBranch<Type>(TTree* tree, \
                          const char* name, Type& variable) { \
    tree->Branch(name, \&variable, \setminus
                  (std::string(name) + suffix).c_str()); \
TYPES WITH SUFFIX
#undef X
#undef TYPES_WITH_SUFFIX
template <typename T>
class STree {
  Public:
    STree(const char* name, const char* title)
        : tree (new TTree(name, title))
    \{ init(); \}
    void init();
    void write() const { tree ->Write(); }
    void print() const { tree ->Print(); }
    void fill(T && structure) const {
        std::lock guard<std::mutex> lock {mutex };
        structure = std::move(structure);
        tree ->Fill();
    }
  private:
    template<class S>
    void addBranch(const char* name, S foo) {
        addTreeBranch(tree , name, structure .*foo);
    }
    TTree * tree_;
    mutable T structure ;
    mutable std::mutex mutex ;
};
#endif
```

### The STree – a wrapper around TTree

- Let's first make a function to add branches of the correct type by matching the type of the variable
- Let's create a wrapper around TTree which provides a fill function that does not depend on the external state but takes the branch variables as inputs
- To use this class, one has to define the tree structure as a struct in C++. Type checking of branches is done at compile time and we avert bugs
- Put lock\_guard in write function to make possible multithreaded filling

```
#ifndef EVENT TREE
#define EVENT TREE
#include "Stree.h"
struct EventRecord {
    unsigned int
                       nElectrons:
    std::vector<float> etaVec:
    std::vector<float> phiVec;
    std::vector<float> ptVec;
};
template<>
void STree<EventRecord>::init() {
    addBranch("n"
                     , &EventRecord::nElectrons);
    addBranch("eta"
                     , &EventRecord::etaVec);
    addBranch("phi"
                     , &EventRecord::phiVec);
    addBranch("pt"
                     , &EventRecord::ptVec);
}
```

```
using EventTree = STree<EventRecord>;
```

# endif

# Defining the tree structure

- Next, we have to define the tree structure and corresponding STree<T>::init() function for this structure
  - You might of course also write a preprocessor macro that does both at once



```
#include "TFile.h"
#include "TRandom.h"
#include "EventTree.h"
```

#include <vector> #include <cmath>

```
class MySimulator {
```

public:

```
MySimulator() : tree ("tree", "tree") {}
    ~MySimulator() { tree .write(); }
    void simulate(int iEvent) const
        // how many electrons?
        const unsigned int nElectrons {4 + rng .Integer(4)};
        std::vector<float> etaVec(nElectrons);
        std::vector<float> phiVec(nElectrons);
        std::vector<float> ptVec(nElectrons);
        for (auto& x : etaVec) x = rng .Uniform() * 5.0 - 2.5;
        for (auto& x : phiVec) x = rng .Uniform() * 2.*M PI - M PI;
        for (auto& x : ptVec) x = rng_.Exp(10.) + 20.;
        tree .fill({ .nElectrons = nElectrons,
                     .etaVec
                                = etaVec,
                      .phiVec
                                 = phiVec,
                      .ptVec
                                 = ptVec
                                               });
   }
  private:
    EventTree tree ;
    mutable TRandom rng ;
};
int main()
    const int nEvents {10000}; // number of events to simulate
   TFile file {"data.root", "RECREATE"};
```

MySimulator simulator{};

```
for(int iEvent = 0; iEvent < nEvents; ++iEvent)</pre>
    simulator.simulate(iEvent);
```

ł

## Our final FP solution

- **Impure** part now mostly **moved** • out to other files
- **fill** function now much more • explicit as it was before
- Code is easier to read now: a • look at the fill function call is enough to get the general idea
- No type-related bugs possible • with the STree wrapper
- On the downside about 10 % slower
- Speed loss can be mitigated by simulating in **multiple threads**, which would not have been possible before

```
#include "TFile.h"
#include "TRandom.h"
#include "EventTree.h"
```

#include <vector> #include <cmath>

```
class MySimulator {
```

public:

ł

```
MySimulator() : tree ("tree", "tree") {}
    ~MySimulator() { tree .write(); }
    void simulate(int iEvent) const
        // how many electrons?
        const unsigned int nElectrons {4 + rng .Integer(4)};
        std::vector<float> etaVec(nElectrons);
        std::vector<float> phiVec(nElectrons);
        std::vector<float> ptVec(nElectrons);
        for (auto& x : etaVec) x = rng .Uniform() * 5.0 - 2.5;
        for (auto& x : phiVec) x = rng .Uniform() * 2.*M PI - M PI;
        for (auto& x : ptVec) x = rng .Exp(10.) + 20.;
        tree .fill({ .nElectrons = nElectrons,
                     .etaVec
                                = etaVec.
                      .phiVec
                                 = phiVec,
                      .ptVec
                                 = ptVec
                                                }):
    }
  private:
    EventTree tree ;
    mutable TRandom rng ;
};
int main()
    const int nEvents {10000}; // number of events to simulate
    const int nThreads {8};
    TFile file {"data.root", "RECREATE"};
    MySimulator simulator{};
    auto simulateEvents = [\&]() {
        for(int iEvent = 0; iEvent < nEvents/nThreads; ++iEvent)</pre>
            simulator.simulate(iEvent);
    };
    std::vector<std::thread> threads:
    for(int i = 0; i < nThreads; ++i) threads.emplace back(simulateEvents);</pre>
    for(auto& thread : threads) thread.join();
```

## Our final FP solution (multi-threaded)

- Easy to speed up your analysis • by factor n if you kept your use of mutable data under control
- **Downside:** the order of events is • now not deterministic anymore, which makes the **random** numbers not deterministic



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import uproot import numpy as np import matplotlib.pyplot as plt import pandas as pd

```
df = uproot.open("data.root")["tree"].pandas.df()
```

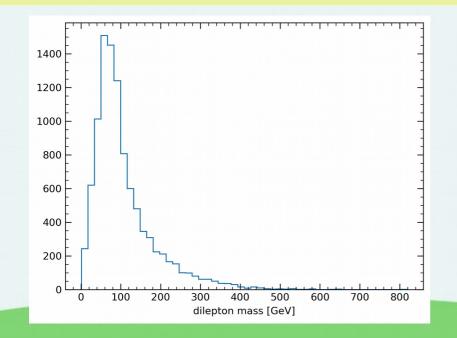
```
def get_mass(leptons):
    leps = leptons.sort_values("pt")[-2:]
```

```
pt_prod = leps["pt"].prod()
eta_diff = leps["eta"].diff().iloc[1]
phi_diff = leps["phi"].diff().iloc[1]
```

return np.sqrt(2\*pt\_prod\*(np.cosh(eta\_diff)-np.cos(phi\_diff)))

```
mass = df.groupby("entry").apply(get_mass).values
```

```
plt.figure()
plt.hist(mass, bins=50, histtype='step')
plt.xlabel("dilepton mass [GeV]")
plt.savefig("dilepton_mass.png")
```



## Making Plot with Python

- The **split-apply-combine** idiom is a nice example of how higheroder functions can abstact away event loops
- In Pandas:
  - Split step is done by df.groupby
  - apply is the higher-order function
  - Combination step is implicit in apply
- **Pandas** provides a very high level of abstaction and is fast enough for datasets with thousands of rows
- Pandas disencourages mutable data, but you can still do it with the inplace=True argument

## Conclusions

- Functional programming is the natural **next level of abstraction** in the evolution of programming languages
- Functional programming languages like Haskell or OCaml (while not being mainstream yet) have significantly inspired the mainstream languages to implement features originating from FP
- These new aspects of C++ and Python can help Physicists write **more bug-free and readable code**, saving time for everyone
- Three most important lessons:
  - Write pure functions
  - The compiler and type-system is your friend
  - Avoid mutating variables if possible
  - Use higher-level functions
- Sticking to these rules results in code that can be easily parallelized